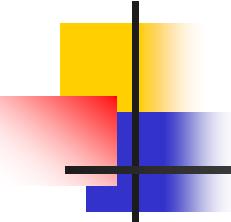


Arithmetic logic and Verilog

CMPE 415

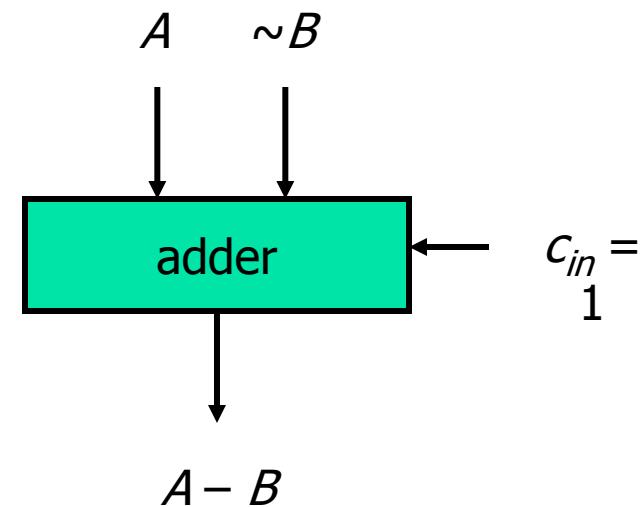
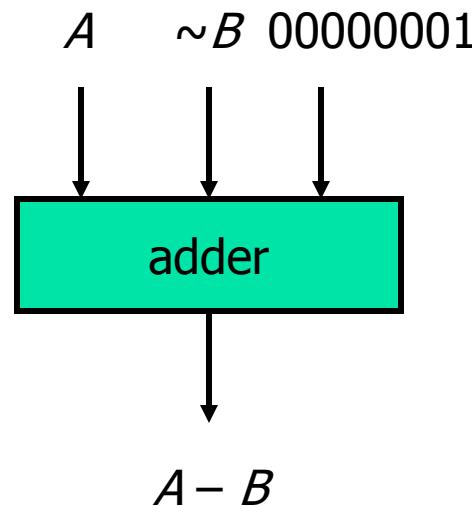


Subtraction

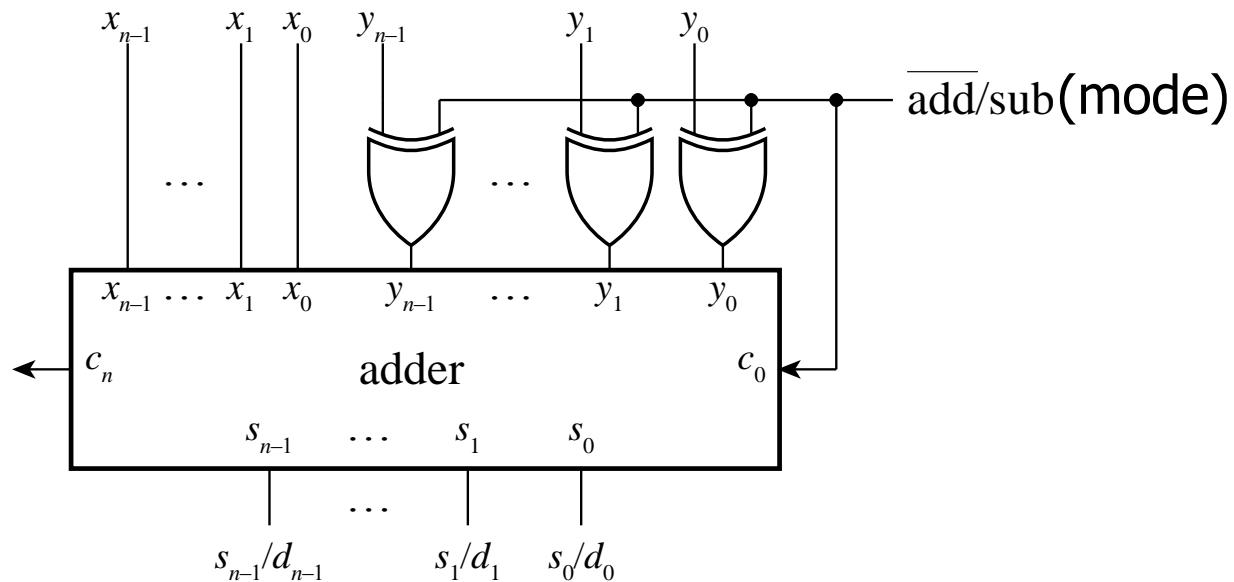
- Essentially the same hardware as an adder
- $A - B = A + (-B)$
- Recall that for 2's complement numbers,
 $-B = (\sim B) + 1$
- So now we have
 $A - B = A + (\sim B) + 1$

Subtraction

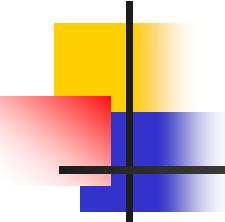
- Often easy to find a place to add in a “1” in the lsb position



Adder/Subtractor Circuits



- Adder can be any of those we have seen
 - depends on constraints



Subtraction in Verilog

```
module adder_subtracter ( output [12:0] s,
                           input  [11:0] x, y,
                           input          mode );
  assign {s} = !mode ? (x + y) : (x - y);
endmodule
```

- When mode=0 => adds
 - else subtracts

Unary Reduction Operator

```
module and4_rtl(y_out, x_in);  
  input [3:0] x_in;  
  output y_out;  
  
  assign y_out = & x_in;  
  
endmodule
```

Equivalents:

```
assign y_out =  
x_in[4] & x_in[2] &  
x_in[1] & x_in[0];
```

```
and(y_out, x_in[4],  
x_in[2], x_in[1],  
x_in[0]);
```

Defining Port Connections

```
module Add_half (sum, c_out, a, b); Module de
```

...

```
half_adder m1(w1, w2, a, b); Ordered conn  
half_adder m2(sum, w3, w1, c_in); Order define  
Module defin
```

```
half_adder m1(.sum(w1), .c_out(w2),  
              .a(a), .b(b));
```

```
half_adder m2(.sum(sum), .c_out(w3),  
              .a(w1), .b(c_in));
```

Ports
Expl
name

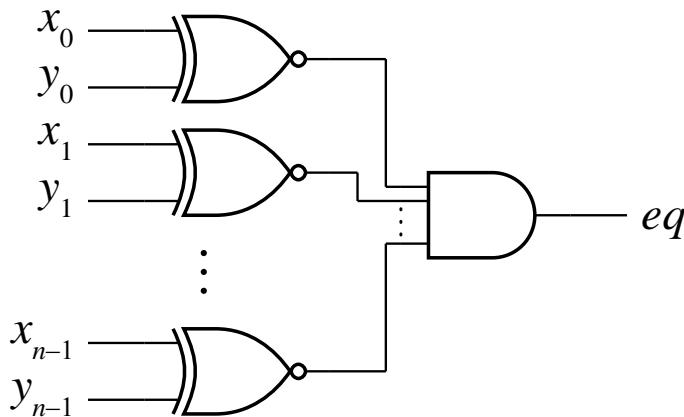
```
half_adder m1(.a(a), .b(b),  
              .sum(w1), .c_out(w2));
```

```
half_adder m2(.sum(sum), .c_out(w3));
```

so ord
doesn'
matter

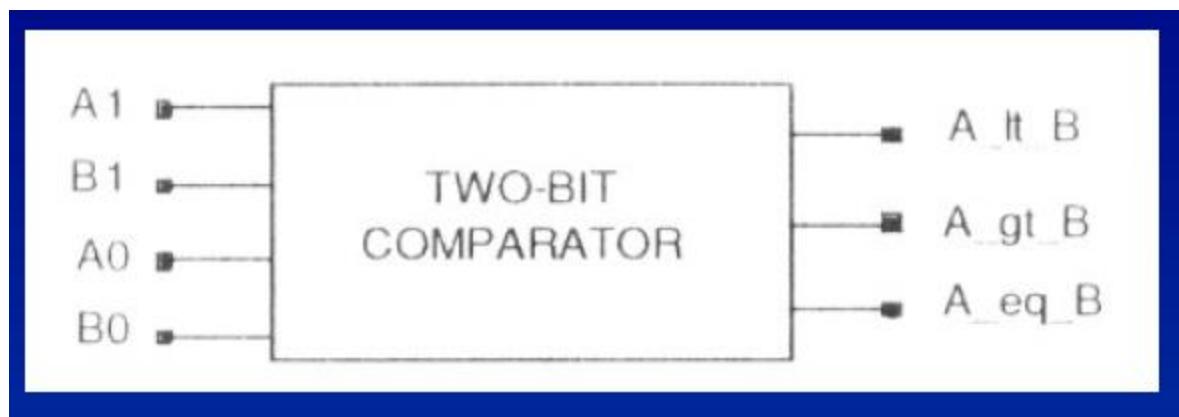
Equality Comparison

- XNOR gate: equality of two bits
 - Apply bitwise to two unsigned numbers



- In Verilog, $x == y$ gives a bit result
 - 1'b0 for false, 1'b1 for true

```
assign eq = x == y;
```



Long

```
module compare_2_str (A_lt_B, A_gt_B,  
A_eq_B, AO, A1, BO, B1);  
    input AO, A1.B0, B1;  
    output A_lt_B, A^gLB, A_eq_B;  
    wire w1, w2, w3, w4, w5, w6, w7;  
    or (A_lt_B, w1, w2, w3);  
    nor (A_gt_B, A_lt_B, A_eq_B);  
    and (A.. eq_B, w4, w5);  
    and (w1, w6, B1);  
    and (w2, w6, w7, BO);  
    and (w3, w7, BO, B1);  
    not (w6, A1);  
    not (w7, AO);  
    xnor (w4, A1.B1);  
    xnor (w5, AO, BO);  
endmodule
```

```
module compare_2a (A_lt_B, A_gt_B, A_eq_B, A1, AO,  
B1, BO);  
    input A1, AO, B1, BO;  
    output A_lt_B, A_gt_B, A..eq_B;  
    assign A_lt_B = (~A1) & B1 | (~A1) & (~AO) & BO  
    | (~AO) & B1 & BO;  
    assign A_gt_B = A1 & (~B1) | AO & (~B1) & (~BO)  
    | A1 & AO & (~BO);  
    assign A_eq_B = (~A1) & (~AO) & (~B1) & (~BO) |  
    (~A1) & AO & (~B1) & BO | A1 | AO & B1 & BO | A1 &  
    (~AO) & B1 & (~BO);  
endmodule
```

Shorter

```
module compare_2b (A_lt_B, A_gt_B, A_eq_B, A1,  
A0, B1, B0);  
input A1,A0,B1,B0;  
output A_lt_B, A_gt_B, A_eq_B;  
assign A_lt_B = ({A1 ,A0} < {B1 ,B0});  
assign A_gt_B = ({A1 ,A0} > {B1 ,B0});  
assign A_eq_B = ({A1 ,A0} == {B1 ,B0});  
endmodule
```

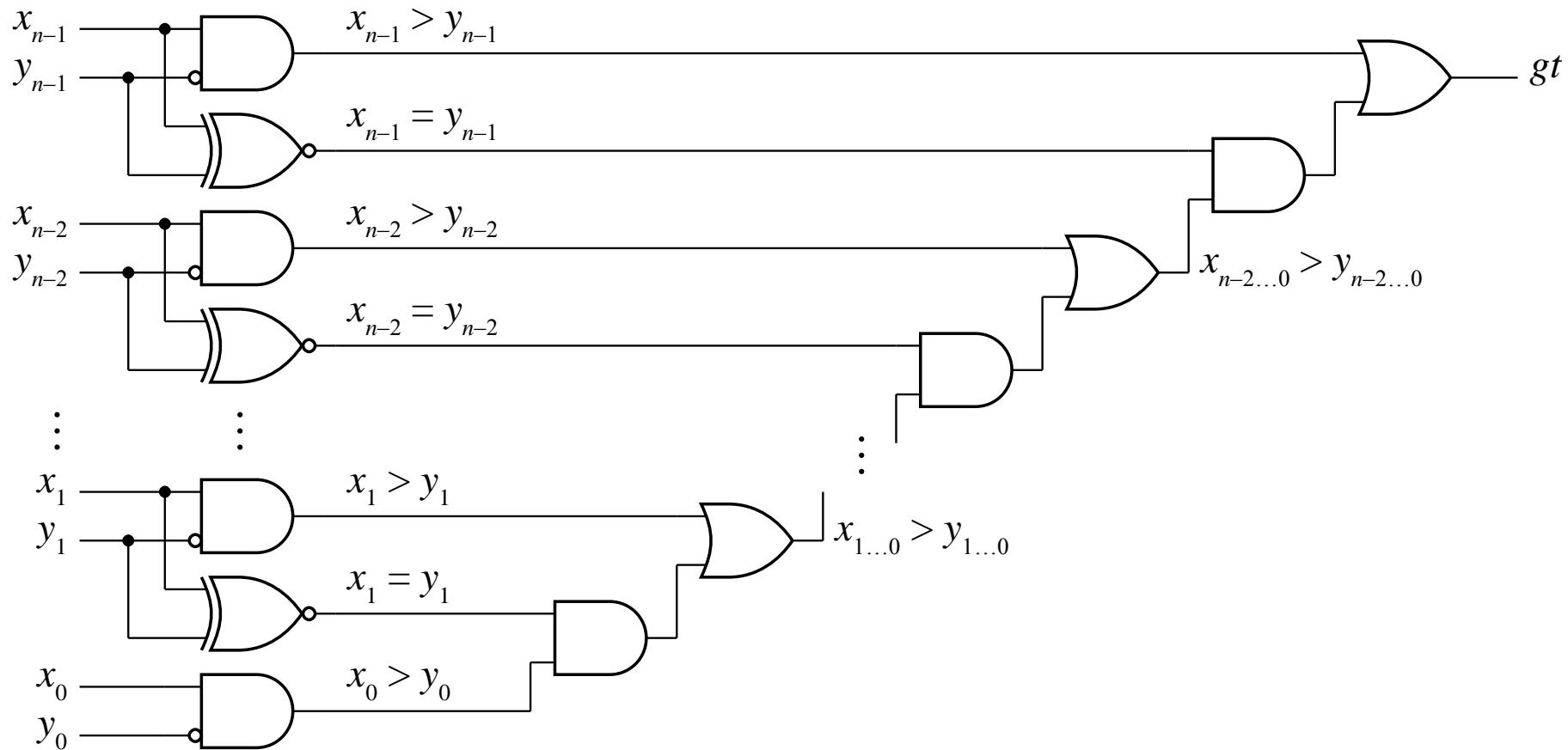
```
module compare_2_ca (A_lt_B, A_gt_B, A_eq_B, A, B);  
input [1:0] A, B;  
output A_lt_B, A_gt_B, A_eq_B;  
assign A_lt_B = (A < B);  
assign A_gt_B = (A > B);  
assign A_eq_B = (A == B);  
endmodule
```

"Algorithmic"

```
module compare_2_algo (A_lt_B, A_gt_B,  
                      A_eq_B, A, B);  
    input [1:0] A, B;  
    output A_lt_B, A_gt_B, A_eq_B;  
    reg A_lt_B, A_gt_B, A_eq_B;  
    always @ (A or B) // Behavior and event expression  
begin  
    A_lt_B = 0;  
    A_gt_B = 0;  
    A_eq_B = 0;  
    if (A==B) A_eq_B = 1;  
    else if (A > B) A_gt_B = 1;  
    else A_lt_B = 1;  
end  
endmodule
```

Inequality Comparison

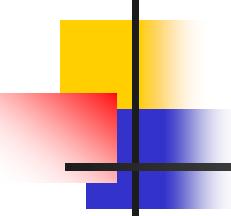
- Magnitude comparator for $x > y$



Comparison Example in Verilog

- Thermostat with target temperature
 - Heater or cooler on when actual temperature is more than 5° from target

```
module thermostat ( output      heater_on, cooler_on,  
                     input [7:0] target, actual );  
    assign heater_on = actual < target - 5;  
    assign cooler_on = actual > target + 5;  
endmodule
```



Scaling by Power of 2

$$x = x_{n-1} 2^{n-1} + x_{n-2} 2^{n-2} + \cdots + x_0 2^0$$

$$2^k x = x_{n-1} 2^{k+n-1} + x_{n-2} 2^{k+n-2} + \cdots + x_0 2^k + (0)2^{k-1} + \cdots + (0)2^0$$

- This is x shifted left k places, with k bits of 0 added on the right
 - *logical shift left* by k places
 - e.g., $00010110_2 \times 2^3 = 00010110000_2$
- Truncate if result must fit in n bits
 - overflow if any truncated bit is not 0

Scaling by Power of 2

$$x = x_{n-1} 2^{n-1} + x_{n-2} 2^{n-2} + \cdots + x_0 2^0$$

$$x / 2^k = x_{n-1} 2^{n-1-k} + x_{n-2} 2^{n-2-k} + \cdots + x_k 2^0 + x_{k-1} 2^{-1} + \cdots + x_0 2^{-k}$$

- This is x shifted right k places, with k bits truncated on the right
 - *logical shift right* by k places
 - e.g., $01110110_2 / 2^3 = 01110_2$
- Fill on the left with k bits of 0 if result must fit in n bits

Scaling in Verilog

- Shift-left (<<) and shift-right (>>) operations
 - result is same size as operand

$$s = 00010011_2 = 19_{10}$$



```
assign y = s << 2;
```

$$y = 01001100_2 = 76_{10}$$



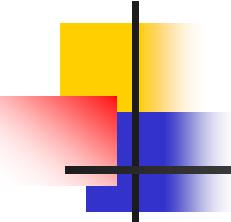
$$s = 00010011_2 = 19_{10}$$



```
assign y = s >> 2;
```

$$y = 000100_2 = 4_{10}$$



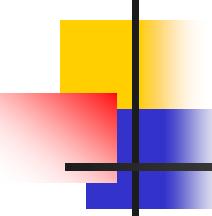


Multipliers

- Widely used in DSP processors, less so in general-purpose processors
- Hardware is typically done the same as you would do it with paper and pencil
- Partial products are copies of the *multiplicand* AND'd by bits of the *multiplier*

<draw picture in notes>

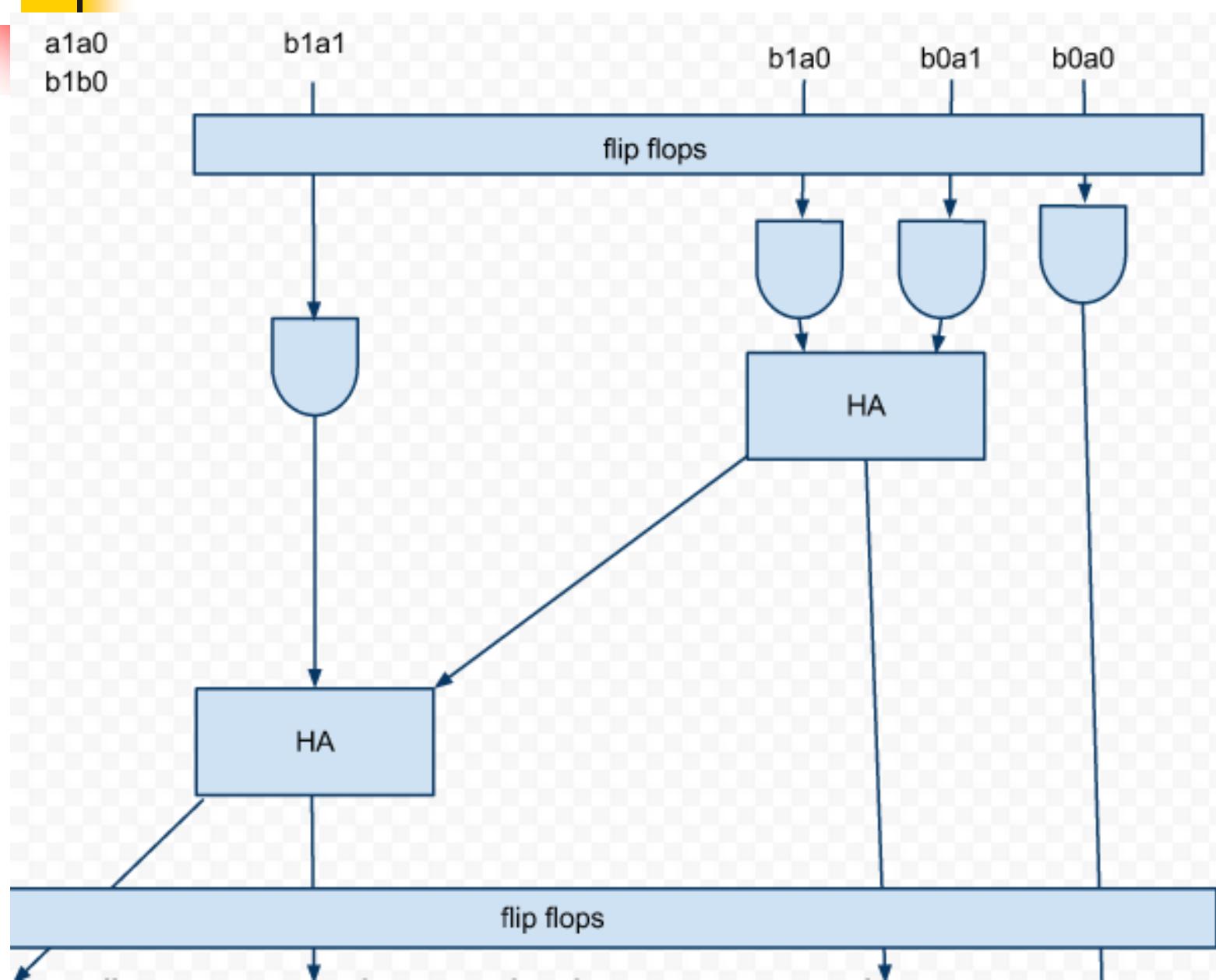
B. Baas, © 2011



Multipliers 3 Main Steps

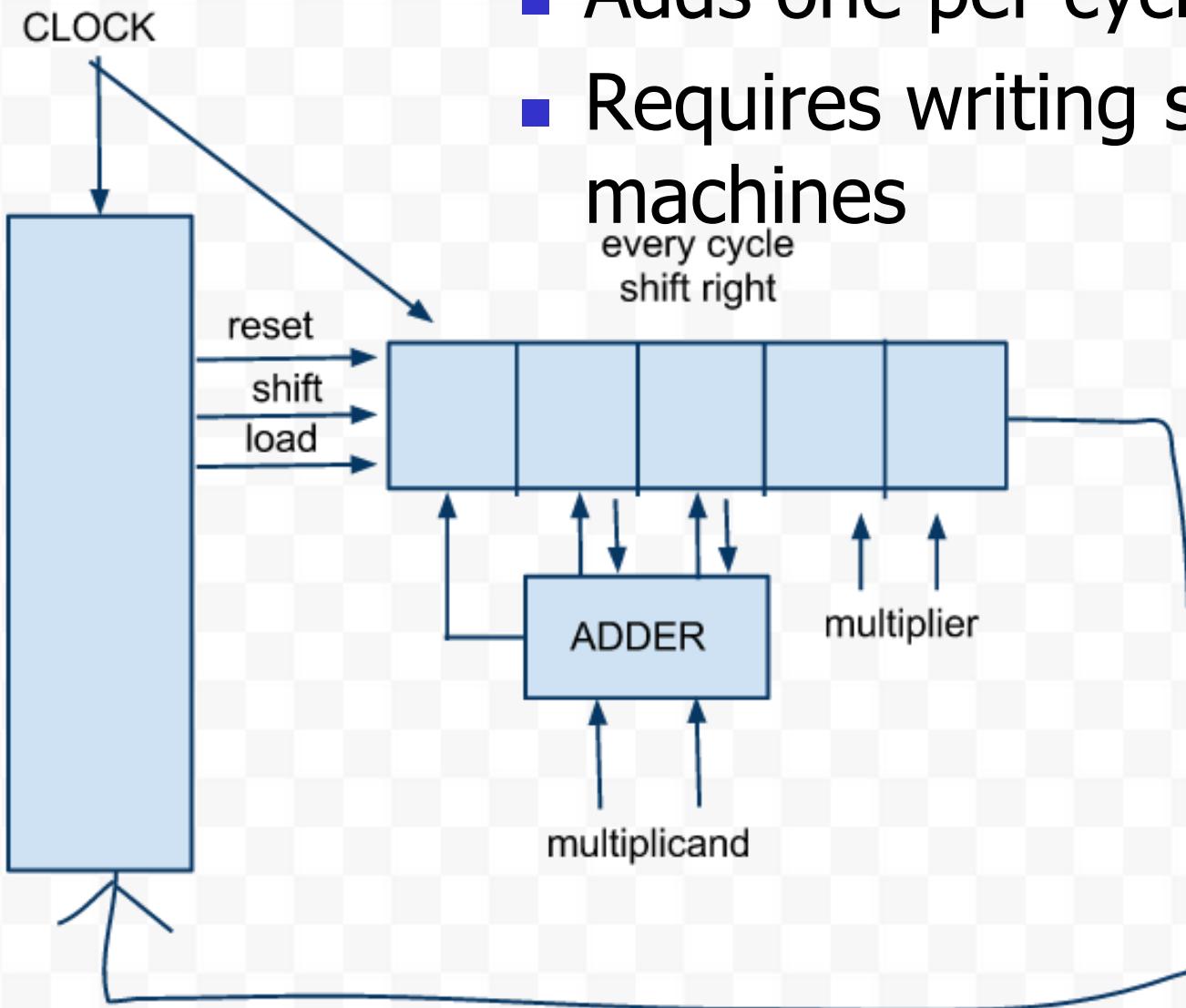
- Generate partial products
- Reduce the partial product array (normally using carry-save addition)
 - Linear array addition
 - Tree addition (Wallace tree)
- Final adder
 - Convert carry-save form to single word form
 - Whichever one you like, probably a faster one
 - Carry-propagate adder (CPA)

Array Multiplier



Serial Multiplier

- Adds one per cycle
- Requires writing state machines
 - every cycle shift right



Unsigned Multiplication

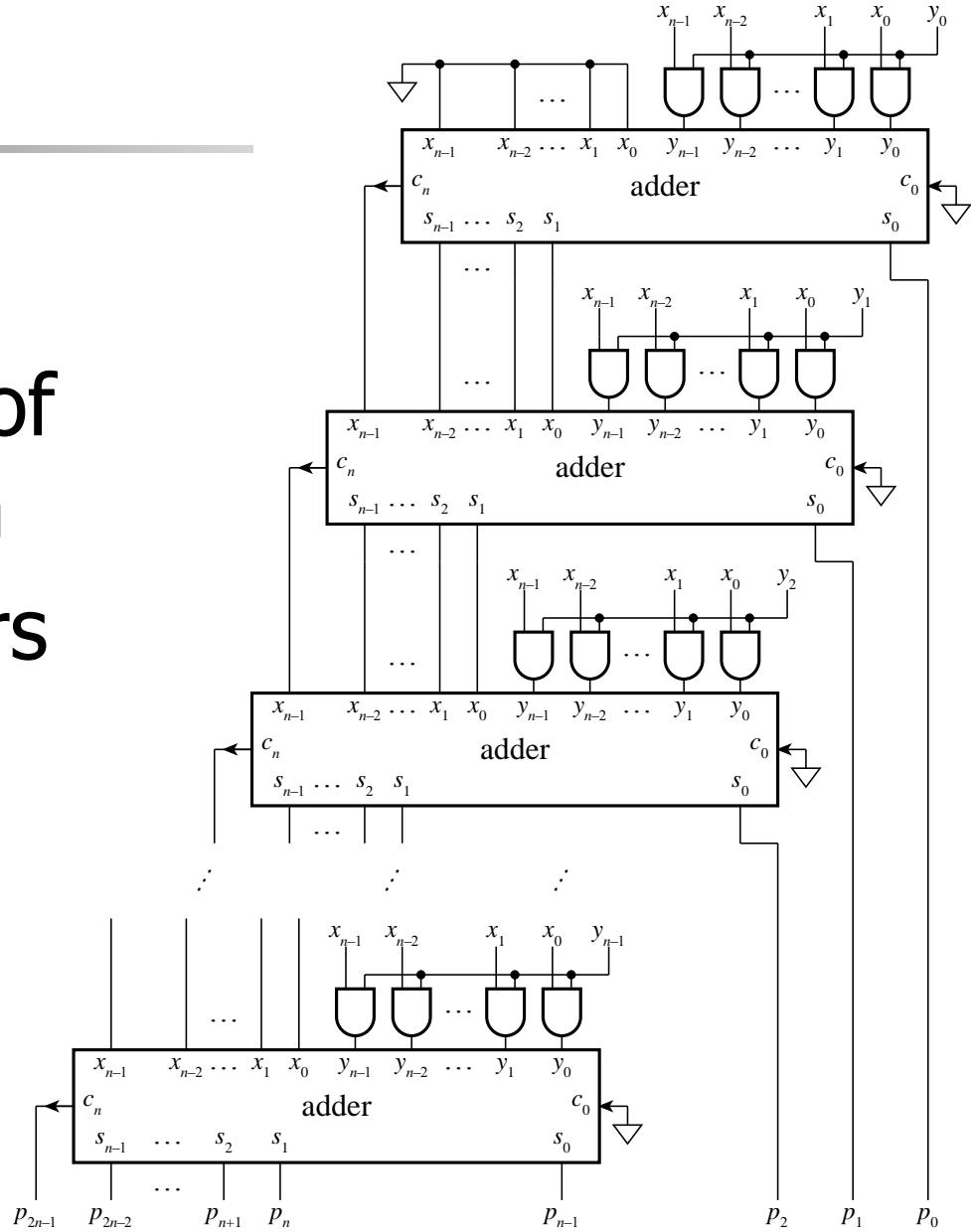
$$xy = x(y_{n-1}2^{n-1} + y_{n-2}2^{n-2} + \dots + y_02^0)$$

$$= y_{n-1}x2^{n-1} + y_{n-2}x2^{n-2} + \dots + y_0x2^0$$

- $y_i x 2^i$ is called a partial product
 - if $y_i = 0$, then $y_i x 2^i = 0$
 - if $y_i = 1$, then $y_i x 2^i$ is x shifted left by i
- Combinational array multiplier
 - AND gates form partial products
 - adders form full product

Unsigned Multiplication

- Adders can be any of those we have seen
- Optimized multipliers combine parts of adjacent adders



Product Size

- Greatest result for n -bit operands:

$$(2^n - 1)(2^n - 1) = 2^{2n} - 2^n - 2^n + 1 = 2^{2n} - (2^{n+1} - 1)$$

- Requires $2n$ bits to avoid overflow
- Multiplying n -bit and m -bit operands
 - requires $n + m$ bits

```
wire [ 7:0] x; wire [13:0] y; wire [21:0] p;
```

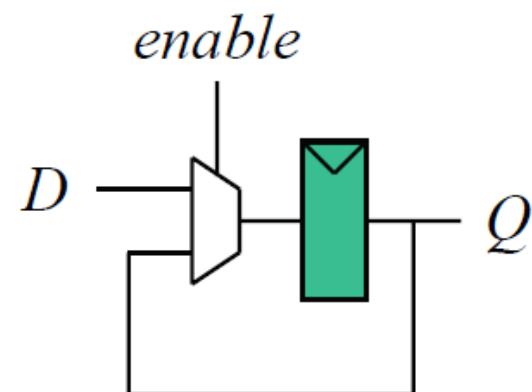
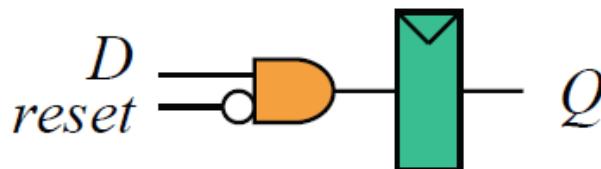
```
...
```

```
assign p = {14'b0, x} * {8'b0, y};
```

```
assign p = x * y; // implicit resizing
```

Reset-able and Enable-able Registers

- Sometimes it is convenient or necessary to have flip-flops with special inputs like *reset* and *enable*
- When designing flip-flops/registers, it is ok (possibly required) for there to be cases where the `always` block is entered, but the `reg` is not assigned
- No fancy code, just make it work
- Normally use synchronous reset instead of asynchronous reset (easier to test)



Reset-able and Enable-able Registers

- Example FF with reset and enable (reset has priority)

```
always @ (posedge clk) begin
    if (reset)           // highest priority
        out <= #1 1'b0;
    else if (enable)
        out <= #1 c_out;
    // ok if no assignment (out holds value)
end
```

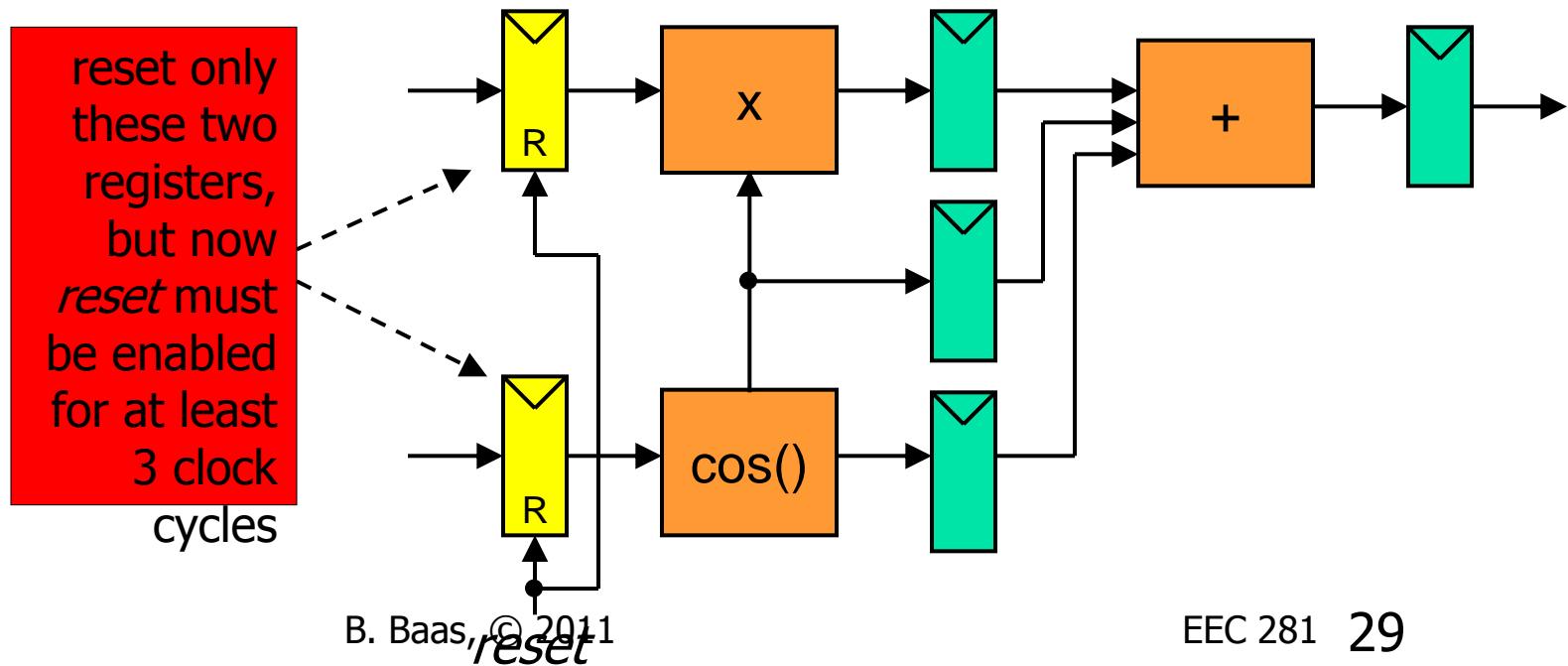
Reset-able and Enable-able Registers

- Example FF with reset and enable (enable has priority)

```
always @ (posedge clk) begin
    if (enable) begin // highest priority
        if (reset)
            out <= #1 1'b0;
        else
            out <= #1 c_out;
    end
    // ok if no assignment (out holds value)
end
```

Reset-able and Enable-able Registers

- Use reset-able FFs only where needed
 - FFs are a little larger and higher power
 - Requires the global routing of the high-fanout *reset* signal



Three types of “case” statements in verilog

1) case

- Normal case statement

2) casez

- Allows use of wildcard “?” character for don’t cares.

```
casez(in)
  4'b1????: out = a;
  4'b01????: out = b;
  4'b00????: out = c;
  default: out = d;
endcase
```

3) casex

- Don’t use it. Could use “z” or “x” logic.
- default
 - Normally set output to an easily-recognizable value (such as x’s) in a default statement to make mistakes easier to spot

Hardwired Complex Functions

- Complex or “arbitrary” functions are not uncommon
 - Examples
 - sin/cos
 - tangent⁻¹
 - log
- theta* —————→ **sin/cos** —————→ *out_real*
————→ *out_imag*

Hardwired Function in Verilog using a Lookup Table

- ```
always @(input) begin
 case (input)
 4'b0000: begin real=3'b100; imag=3'b001; end
 4'b0001: begin real=3'b000; imag=3'b101; end
 4'b0010: begin real=3'b110; imag=3'b011; end
 ...
 default: begin real=3'bxxx; imag=3'bxxx; end
 endcase
end
```
- Often best to write a matlab program to write the verilog table as plain text
  - You will need several versions to get it right
  - Easy to adapt to other specifications
- Not efficient for very large tables
- Tables with data that is less random will have smaller synthesized area