

# **CMSC 611: Advanced Computer Architecture**

## Branch Prediction

# Correlating Predictors

```
If (aa == 2)
    aa = 0;
If (bb == 2)
    bb = 0;
If (aa != bb) {
```



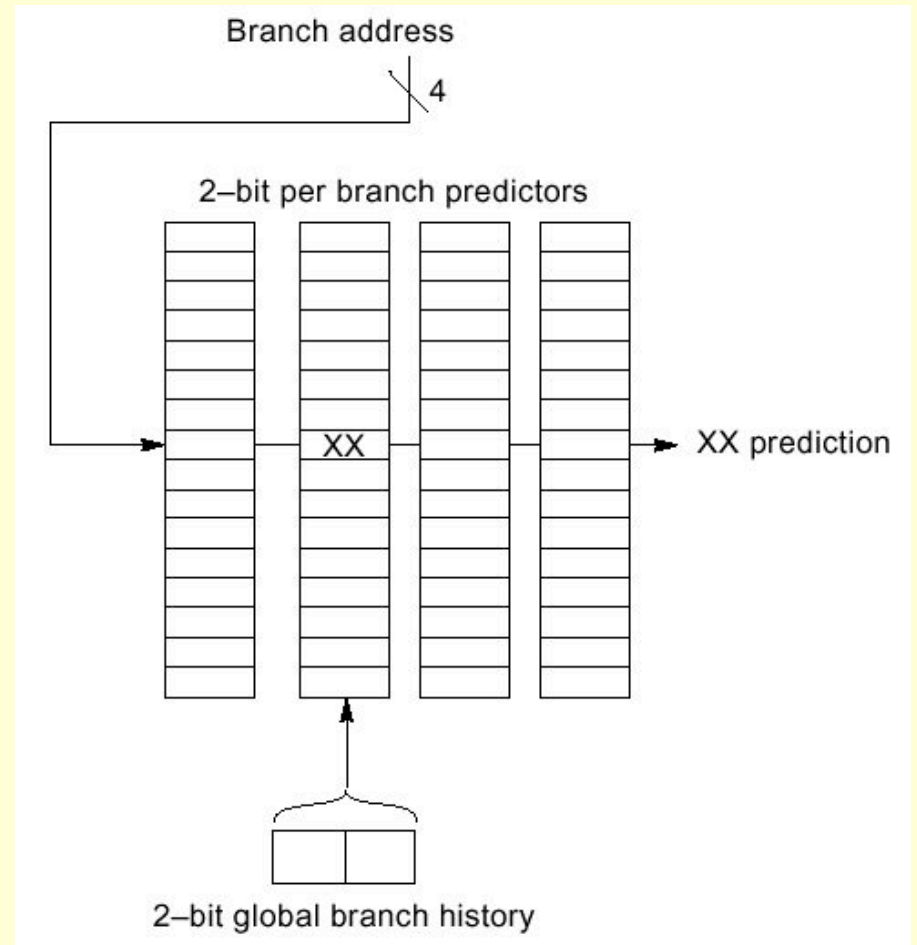
```
DSUBUI    R3, R1, #2
BNEZ      R3, L1      ; branch b1 (aa!=2)
ANDI      R1, R1, #0  ; aa=0
L1: SUBUI  R3, R2, #2
BNEZ      R3, L2      ; branch b2 (bb!=2)
ANDI      R2, R2, #0  ; bb=0
L2: SUBU   R3, R1, R2  ; R3=aa-bb
BEQZ      R3, L3      ; branch b3 (aa==bb)
```

- The behavior of branch b3 is correlated with the behavior of b1 and b2
- Clearly if both branches b1 and b2 are untaken, then b3 will be taken
- A predictor that uses only the behavior of a single branch to predict the outcome of that branch can never capture this behavior
- Branch predictors that use the behavior of other branches to make a prediction are called correlating or two-level predictors

**Hypothesis: recent branches are correlated; that is, behavior of recently executed branches affects prediction of current branch**

# (2,2) Correlating Predictors

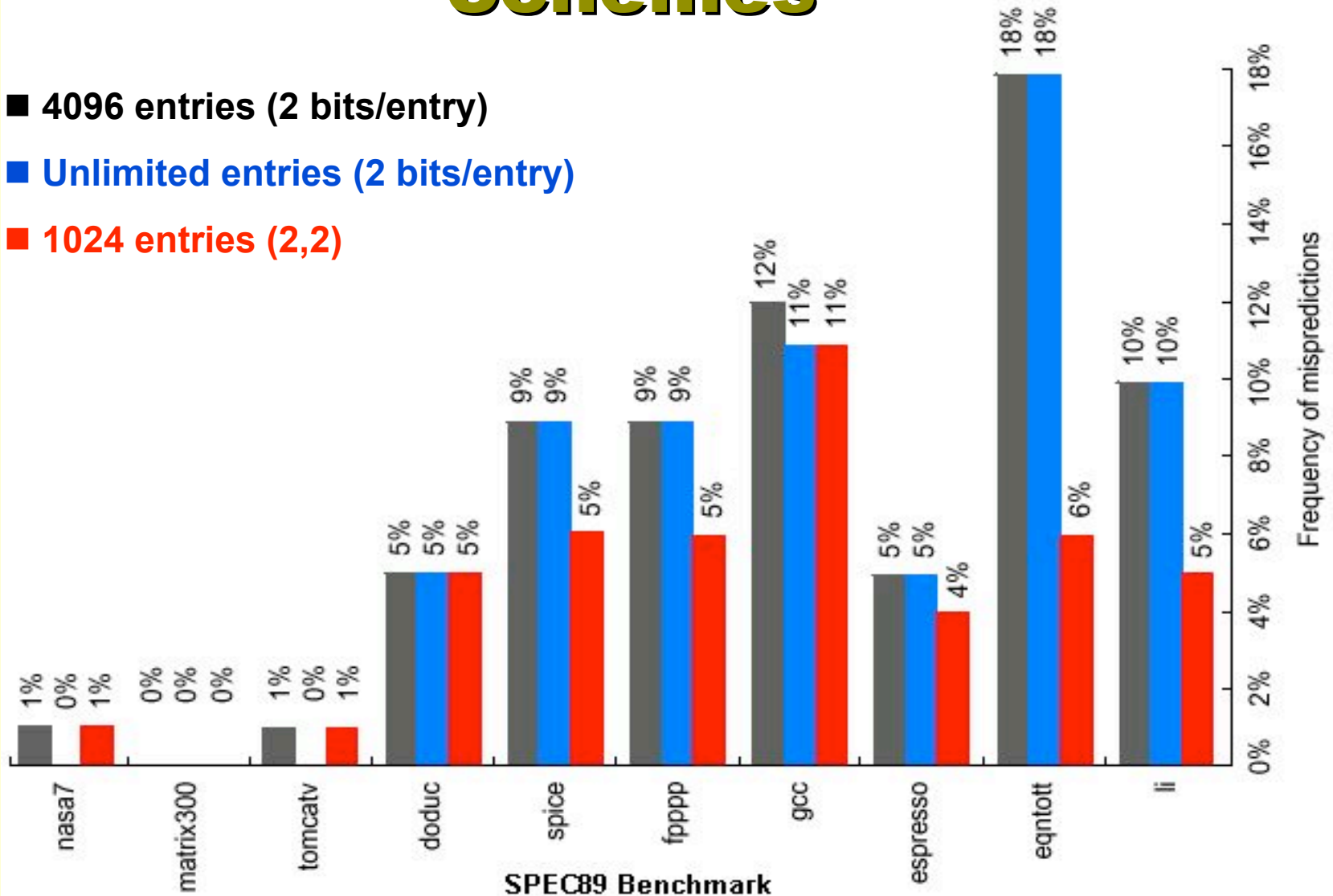
- Record  $m$  most recently executed branches as taken or not taken, and use that pattern to select the proper branch history table
- $(m,n)$  predictor means record last  $m$  branches to select between  $2^m$  history tables each with  $n$ -bit counters
  - Old 2-bit branch history table is a  $(0,2)$  predictor
- In a  $(2,2)$  predictor, the behavior of recent branches selects between, four predictions of next branch, updating just that prediction



Total size =  $2^m \times n \times \#$  prediction entries selected by branch address

# Accuracy of Different Schemes

- 4096 entries (2 bits/entry)
- Unlimited entries (2 bits/entry)
- 1024 entries (2,2)



# Example

- Assume that d has values 0, 1, or 2 (alternating between 0, 2)
- Assume that the sequence will be executed repeatedly
- Ignore all other branches including those causing the sequence to repeat
- All branches are initially predicted to untaken state

```
if (d==0)           BNEZ      R1, L1           ; branch b1 (d!=0)
    d=1;           DADDI     R1, R0, #1       ; d==0, sp d=1
if (d==1)           L1: DSUBUI  R3, R1, #1
                    BNEZ     R3, L2           ; branch b2 (d!=1)
                    ....
                    L2:
```

# Example

## With a single bit predictor

d=?	b1 prediction	b1 action	New b1 prediction	b2 prediction	b2 action	New b2 prediction
2	NT	T	T	NT	T	T
0	T	NT	NT	T	NT	NT
2	NT	T	T	NT	T	T
0	T	NT	NT	T	NT	NT

- *All branches are mispredicted*

```
if (d==0)          BNEZ      R1, L1          ; branch b1 (d!=0)
    d=1;           DADDI      R1, R0, #1      ; d==0, sp d=1
if (d==1)          L1: DSUBUI   R3, R1, #1
                   BNEZ      R3, L2          ; branch b2 (d!=1)
                   .....
                   L2:
```

# Example

## With one bit predictor with one bit of correlation

d=?	b1 prediction	b1 action	New b1 prediction	b2 prediction	b2 action	New b2 prediction
2	NT/NT	T	T/NT	NT/NT	T	NT/T
0	T/NT	NT	T/NT	NT/T	NT	NT/T
2	T/NT	T	T/NT	NT/T	T	NT/T
0	T/NT	NT	T/NT	NT/T	NT	NT/T

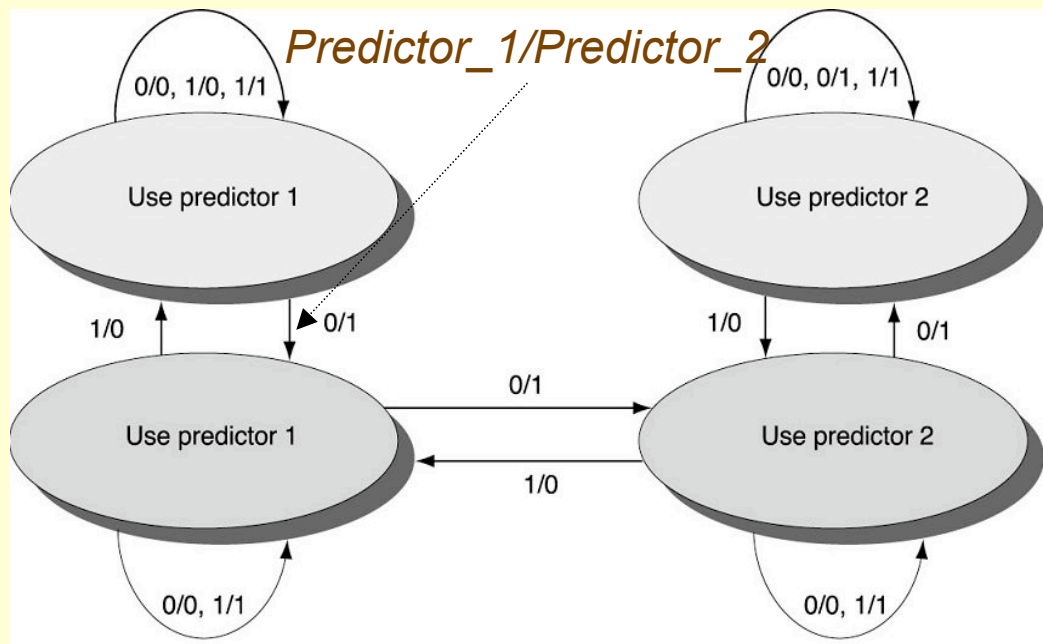
- *Except for first iteration, all branches are correctly predicted*

```

if (d==0)
    d=1;
if (d==1)
    BNEZ    R1, L1          ; branch b1 (d!=0)
            DADDI   R1, R0, #1      ; d==0, sp d=1
L1: DSUBUI  R3, R1, #1
            BNEZ    R3, L2          ; branch b2 (d!=1)
            ....
L2:
    
```

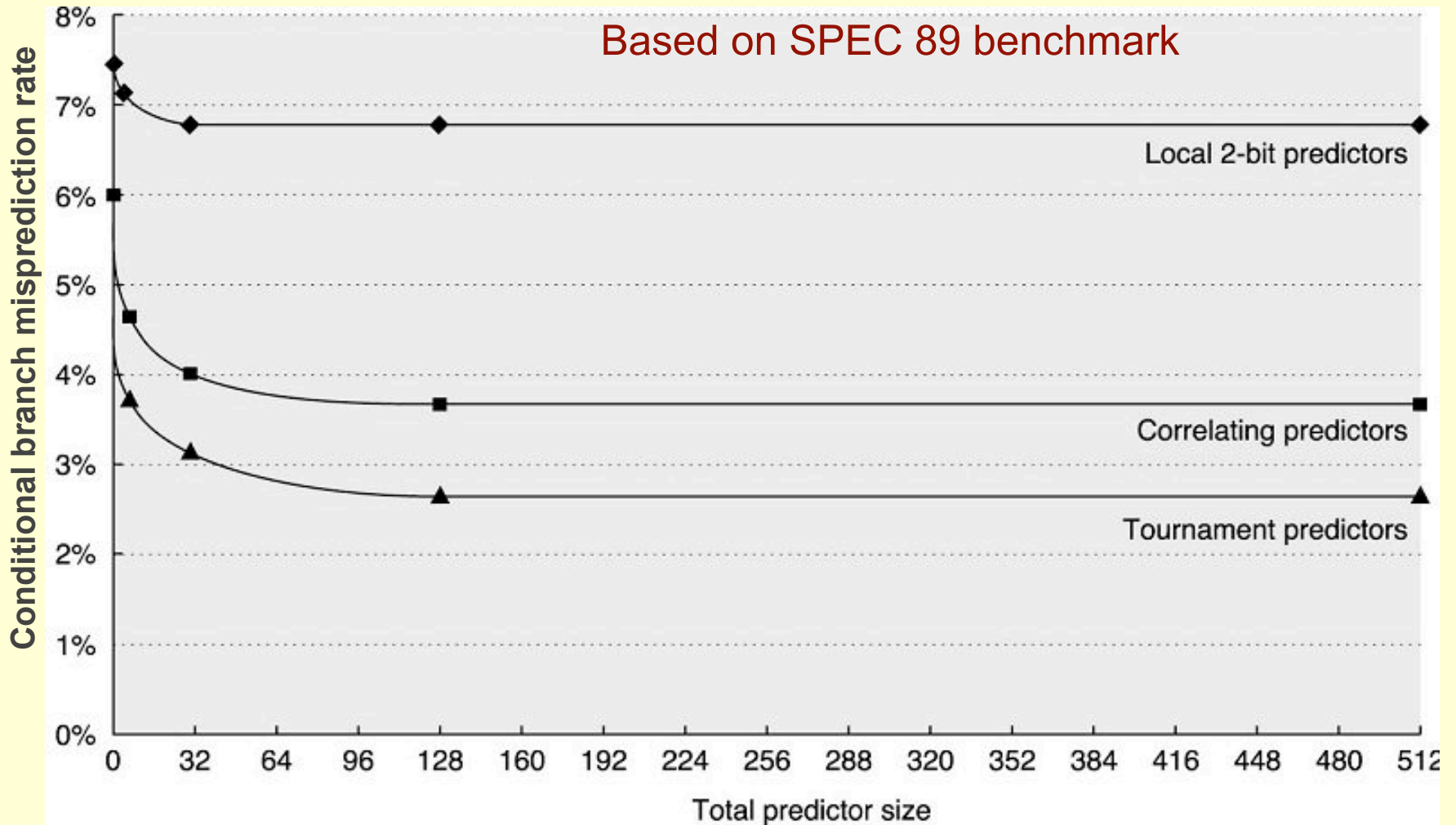
# Tournament Predictors

- Multilevel branch predictors use several levels of branch prediction tables together with an algorithm to choose among them
- Tournament selectors are the most popular form of multilevel branch predictors (e.g. DEC Alpha 21264)
- Tournament predictors combines both local and global predictor
- Selection between the two predictors are based on a selector (2-bit counter)
- Make a transition with two wrong prediction using the current table for which the correct prediction would have been possible using the other predictor





# Performance of Tournament Predictors



Tournament predictors slightly outperform correlating predictors