

Lecture 3

➤ Logic

- Predicates and Quantified Statements
- Statements with Multiple Quantifiers

Reading (Epp's textbook)
3.1 - 3.3

➤ Introduction to Proofs

Reading (Epp's textbook)
4.1-4.2

Propositional Functions

- A **predicate** is a sentence that contains a finite number of variables and becomes a statement when specific values are substituted for the variables.
 - e.g.: $x - 3 > 5$.
- Let us call this propositional function $P(x)$, where P is the **predicate** ($x - 3 > 5$) and x is the **variable**.

What is the truth value of $P(2)$?

false

What is the truth value of $P(8)$?

false

What is the truth value of $P(9)$?

true

Propositional Functions

- Let us consider the propositional function $Q(x, y, z)$ defined as:
 - $x + y = z$.
- Here, Q is the **predicate** and x , y , and z are the **variables**.

What is the truth value of $Q(2, 3, 5)$?

true

What is the truth value of $Q(0, 1, 2)$?

false

What is the truth value of $Q(9, -9, 0)$?

true

Truth Set

- If $P(x)$ is a propositional function and x has domain D
- The **truth set** of $P(x)$ is denoted
 - $\{x \in D \mid P(x)\}$
 - **Note:** $\forall x \in \text{in the truth set of } P(x), P(x)$ should be true.

Universal Quantification

- Let $P(x)$ be a propositional function.
- **Universally quantified sentence:**
 - For all x in the universe of discourse $P(x)$ is true.
- Using the universal quantifier $\forall$:
 - $\forall x \in D, P(x)$
 - **Note:** $\forall x \in D, P(x)$ is either true or false, so it is a statement (proposition), not a propositional function.
- A value for x for which $P(x)$ is false is called a **counterexample** to the universal statement.
 - Statement: “All birds can fly.”
 - Disproved by counterexample: Penguin.

Universal Quantification

➤ Example:

- $S(x)$: x is a UMBC student.
- $G(x)$: x is a genius.

➤ What does $\forall x, (S(x) \rightarrow G(x))$ mean ?

“If x is a UMBC student, then x is a genius.”

or

“All UMBC students are geniuses.”

Existential Quantification

➤ **Existentially quantified sentence:**

- There exists an x in the universe of discourse for which $P(x)$ is true.

➤ Using the existential quantifier $\exists$:

- $\exists x \in D$ such that $P(x)$

“There is an x such that $P(x)$.”

“There is at least one x such that $P(x)$.”

➤ **Note:** $\exists x \in D$ such that $P(x)$ is either true or false, so it is a statement (proposition), but not a propositional function.

Existential Quantification

➤ Example:

- $P(x)$: x is a UMBC professor.
- $G(x)$: x is a genius.

➤ What does $\exists x$ such that $(P(x) \wedge G(x))$ mean ?

“There is an x such that x is a UMBC professor and x is a genius.”

or

“At least one UMBC professor is a genius.”

Negation

➤ Negation of a Universal Statement

$\neg (\forall x \in D, P(x))$ is logically equivalent to
 $\exists x \in D$ such that $\neg P(x)$.

➤ Negation of an Existential Statement

$\neg (\exists x \in D \text{ such that } P(x))$ is logically equivalent to
 $\forall x \in D, \neg P(x)$.

Negation of a Universal Cond. Statement

$\neg (\forall x, \text{if } P(x) \text{ then } Q(x))$ is logically equivalent to
 $\exists x \text{ such that } P(x) \text{ and } \neg Q(x).$

➤ Proof

$\neg (\forall x, \text{if } P(x) \text{ then } Q(x)) \equiv \exists x \text{ such that, } \neg(P(x) \rightarrow Q(x)).$

But

$$\sim(P(x) \rightarrow Q(x)) \equiv P(x) \wedge \sim Q(x)$$

Negation of a Multiply-Quantified Statements

$\sim(\forall x \text{ in } D, \exists y \text{ in } E \text{ such that } P(x, y))$ is logically equivalent to

$$\exists x \text{ in } D \text{ such that } \forall y \text{ in } E, \sim P(x, y).$$

And

$\sim(\exists x \text{ in } D \text{ such that } \forall y \text{ in } E, P(x, y))$ is logically equivalent to

$$\forall x \text{ in } D, \exists y \text{ in } E \text{ such that } \sim P(x, y).$$

Methods of Proof: Some Definitions

- For all objects A , B , and C , (1) $A = A$, (2) if $A = B$ then $B = A$, and (3) if $A = B$ and $B = C$, then $A = C$.
- The set of all integers is closed under addition, subtraction, and multiplication.
- If n is an integer, then
$$n \text{ is even} \Leftrightarrow \exists \text{ an integer } k \text{ such that } n = 2k.$$
$$n \text{ is odd} \Leftrightarrow \exists \text{ an integer } k \text{ such that } n = 2k + 1.$$
- n is prime $\Leftrightarrow \forall$ positive integers r and s , if $n = rs$ then either $r = 1$ and $s = n$ or $r = n$ and $s = 1$.
- n is composite $\Leftrightarrow \exists$ positive integers r and s such that $n = rs$ and $1 < r < n$ and $1 < s < n$.

Proving Existential Statements

- A statement in the form:

$$\exists x \in D \text{ such that } P(x)$$

is true if, and only if,

$P(x)$ is true for at least one x in D .

- Example

- Prove the following: $\exists$ an even integer n that can be written in two ways as a sum of two prime numbers.
- Solution: Let $n = 10$. Then $10 = 5 + 5 = 3 + 7$ and 3, 5, and 7 are all prime numbers.

- This method is called **constructive proof of existence!**

Proving Universal Statements

To prove a statement of the form

$$\forall x \text{ in } D, \text{ if } P(x) \text{ then } Q(x)$$

- I. When D is finite or when only a finite number of elements satisfy $P(x)$, such a statement can be proved by the **method of exhaustion**.
- II. **Method of Generalizing from the Generic Particular:** suppose x is a *particular* but *arbitrarily chosen* element of the set, and show that x satisfies the property.

When the method of generalizing from the generic particular is applied to a property of the form “If $P(x)$ then $Q(x)$,” the result is the **method of direct proof**.

Disproving a Statement

- To **disprove** a universal statement of the form:

$$\forall x \text{ in } D, \text{ if } P(x) \text{ then } Q(x).$$

we should show that the negation of the statement is true

$$\exists x \text{ in } D \text{ such that } P(x) \text{ and not } Q(x).$$

- Such an x is called a **counterexample**.

- To **disprove** an existential statement of the form:

$$\exists x \text{ in } D \text{ such that } P(x).$$

we should show that the negation of the statement is true

$$\forall x \text{ in } D, \neg P(x).$$

- Use the **Method of Generalizing from the Generic Particular**.

Disproving a Statement: Example

➤ Disprove the universal statement:

$$\forall a, b \in R, \text{ if } a^2 = b^2 \text{ then } a = b.$$

Solution: We should show that the negation of the statement is true.

$$\exists a, b \in R \text{ such that } a^2 = b^2 \text{ and not } a = b.$$

Counterexample: Let $a = 1$ and $b = -1$.

Disproving a Statement: Example

➤ Disprove the existential statement:

$\exists x \text{ in } \mathbb{Z}^+ \text{ such that } x^2 + 3x + 2 \text{ is prime.}$

Solution: We should show that the negation of the statement is true

$\forall x \in \mathbb{Z}^+, x^2 + 3x + 2$ **is not** prime (composite).

Method of Generalizing from the Generic Particular:

Suppose x is any (particular but arbitrarily chosen) positive integer. Then $x^2 + 3x + 2 = (x + 1)(x + 2)$. Both $(x + 1)$ and $(x + 2)$ are integers and are > 1 (since $x \geq 1$). Thus $x^2 + 3x + 2$ is a product of two integers each greater than 1, and so $x^2 + 3x + 2$ is not prime.

Method of Direct Proof

A statement in the form $\forall x \text{ in } D, p \rightarrow q$ can be proved by showing that if p is true, then q is also true.

Example: Give a direct proof of the theorem “ $\forall n \text{ in } \mathbb{Z}$, if n is odd, then n^2 is odd.”

Idea: Assume that the hypothesis of this implication is true (n is odd). Then use rules of inference and known theorems to show that q must also be true (n^2 is odd).

Direct Proof Example

- n is odd.
- Then $n = 2k + 1$, where k is an integer.
- Consequently, $n^2 = (2k + 1)^2$.
$$= 4k^2 + 4k + 1$$
$$= 2(2k^2 + 2k) + 1$$
- ✓ Since n^2 can be written in this form, it is odd.

Indirect Proof

A statement in the form $\forall x \text{ in } D, p \rightarrow q$ is equivalent to its **contra-positive** $\neg q \rightarrow \neg p$.

- Therefore, we can prove $p \rightarrow q$ by showing that whenever q is false, then p is also false.

Example:

Give an indirect proof of the theorem
“ $\forall n \text{ in } \mathbb{Z}$, if $3n + 2$ is odd, then n is odd.”

Idea: Assume that the conclusion of this statement is false (n is even). Then use rules of inference and known theorems to show that p must also be false ($3n + 2$ is even).

Indirect Proof Example

- n is even.
- Then $n = 2k$, where k is an integer.
- It follows that
$$\begin{aligned} 3n + 2 &= 3(2k) + 2 \\ &= 6k + 2 \\ &= 2(3k + 1) \end{aligned}$$
- Therefore, $3n + 2$ is even.
- ✓ We have shown that the contrapositive of the statement is true, so the implication itself is also true (If $2n + 3$ is odd, then n is odd).

Rational Numbers

- A real number r is rational $\Leftrightarrow \exists$ integers a and b such that $r = \frac{a}{b}$ and $b \neq 0$.
- Every integer is a rational number.
- The sum of any two rational numbers is rational.
- The double of a rational is rational.