

Lecture 2

➤ Logic

- Compound Statements
- Conditional Statements
- Valid & Invalid Arguments
- Digital Logic Circuits

Reading (Epp's textbook)
2.1 - 2.4

Logic

- Logic is a system based on **statements**.
- A **statement** (or **proposition**) is a sentence that is true or false but not both.
- We say that the **truth value** of a statement is either true (T) or false (F).
- Corresponds to **1** and **0** in digital circuits
- Examples:

- “520 < 499”

Is this a statement?

Yes

What is the truth value?

F

- “ $x + y = 5$ ”

Is this a statement?

No

What is the truth value?

?

- “ $x < y$ if and only if $y > x$ ”

Is this a statement?

Yes

What is the truth value?

T

Compound Statements

- One or more statements can be combined to form a single **compound statement**.
- We formalize this by denoting statements with letters such as **p, q, r, s**, and introducing several **logical operators**.
- We will examine the following logical operators:

Logical Operator	Denote	Symbol
Negation	NOT	$\sim$ or $\neg$
Conjunction	AND	$\wedge$
Disjunction	OR	$\vee$
Exclusive or	XOR	$\oplus$

The **order of operations** can be overridden through the use of parentheses.

- These logical operators are used to build more complicated logical expressions out of simpler ones.

Truth Table

- A **statement form** is an expression made up of statement variables (such as p , q , and r) and logical connectives (such as $\sim$, $\wedge$, and $\vee$).
- A **Truth table** for a given statement form displays the truth values that correspond to all possible combinations of truth values for its component statement variables.

Truth table for NOT

p	$\sim p$
T	F
F	T

Truth table for AND

p	q	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F

Truth table for OR

p	q	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

Exclusive or (XOR)

➤ When or is used in its exclusive sense:

$P \oplus q$ means “p or q **but not** both”

$$p \oplus q = (p \vee q) \wedge \sim(p \wedge q).$$

Truth table for XOR

p	q	$p \vee q$	$p \wedge q$	$\sim(p \wedge q)$	$p \oplus q$
T	T	T	T	F	F
T	F	T	F	T	T
F	T	T	F	T	T
F	F	F	F	T	F

Truth table for $(r \vee (\sim(p \wedge q)))$?

Logical Equivalence

- Two *statement forms* are called **logically equivalent** if, and only if, they have identical truth values for each possible substitution of statements for their statement variables. The logical equivalence of statement forms P and Q is denoted by writing $P \equiv Q$.
- **(Showing equivalence)** Are the statements of the forms $\sim(p \wedge q)$ and $\sim p \vee \sim q$ logically equivalent?

p	q	$\sim p$	$\sim q$	$p \wedge q$	$\sim(p \wedge q)$	$\sim p \vee \sim q$
T	T	F	F	T	F	F
T	F	F	T	F	T	T
F	T	T	F	F	T	T
F	F	T	T	F	T	T

De Morgan's Laws:

$$\sim(p \wedge q) \equiv \sim p \vee \sim q$$

and

$$\sim(p \vee q) \equiv \sim p \wedge \sim q$$

Tautologies and Contradictions

- A **tautology** is a statement form that is always true. A statement whose form is a tautology is a **tautological statement (denoted by t)**.
- A **contradiction** is a statement form that is always false. A statement whose form is a contradiction is a **contradictory statement (denoted by c)**.

p	$\sim p$	$p \vee \sim p$	$p \wedge \sim p$
T	F	T	F
F	T	T	F

Simplifying Statement Forms

$$\sim(q \wedge t) \wedge \underline{\sim(\sim q \wedge p)} \wedge (p \vee q)$$

- By De Morgan's Law $\equiv \sim(q \wedge t) \wedge (\underline{\sim(\sim q)} \vee \sim p) \wedge (p \vee q)$
- By the Double negative law $\equiv \underline{\sim(q \wedge t)} \wedge (q \vee \sim p) \wedge (p \vee q)$
- By the Identity Law $\equiv \sim q \wedge \underline{(q \vee \sim p)} \wedge \underline{(p \vee q)}$
- By the Distributive Law $\equiv \sim q \wedge (q \vee \underline{\sim p \wedge p})$
- By the Negation Law $\equiv \sim q \wedge \underline{(q \vee c)}$
- By the Identity Law $\equiv \underline{\sim q \wedge q}$
- By the Negation Law $\equiv c$

Conditional statements

- **Implication (if - then):** “If p then q ” or “ p implies q ” and is denoted $p \rightarrow q$. It is false when p is true and q is false; otherwise it is true.
- We call p the **hypothesis** of the conditional and q the **conclusion**.

Truth table for $p \rightarrow q$

p	q	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

Construct a truth table for the statement form $p \rightarrow q \equiv \sim p \vee q$.

Conditional statements

Suppose a conditional statement of the form “If p then q ” is given.

- **The Negation** of the conditional statement is logically equivalent to “ p and not q .”

Symbolically: $\sim(p \rightarrow q) \equiv p \wedge \sim q$.

- **The Contrapositive** of the conditional statement is “If $\sim q$ then $\sim p$.”

Symbolically: $p \rightarrow q$ is $\sim q \rightarrow \sim p$.

- **The Converse** of the conditional statement is “If q then p .”

Symbolically: $p \rightarrow q$ is $q \rightarrow p$.

- **The Inverse** of the conditional statement is “If $\sim p$ then $\sim q$.”

Symbolically: $p \rightarrow q$ is $\sim p \rightarrow \sim q$.

Conditional statements

- **Biconditional (if and only if):** “ p if, and only if, q ” and is denoted $p \leftrightarrow q$. It is true if both p and q have the same truth values and is false if p and q have opposite values.

Truth table for $p \leftrightarrow q$

p	q	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

$$p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$$

Logical Equivalence involving $\rightarrow, \leftrightarrow$

- Are the statements of the forms $p \leftrightarrow q$ and $(p \rightarrow q) \wedge (q \rightarrow p)$ logically equivalent?

p	q	$p \rightarrow q$	$q \rightarrow p$	$p \leftrightarrow q$	$(p \rightarrow q) \wedge (q \rightarrow p)$
T	T	T	T	T	T
T	F	F	T	F	F
F	T	T	F	F	F
F	F	T	T	T	T

Truth Table

Order of Logical Operators

1. $\sim$ Evaluate negations first.
2. $\vee, \wedge$ Evaluate $\vee$ and $\wedge$ second. If both are present, parentheses may be needed.
3. $\rightarrow, \leftrightarrow$ Evaluate $\rightarrow$ and $\leftrightarrow$ third. If both are present, parentheses may be needed.

Arguments

- An **argument** is a sequence of statements, and an **argument form** is a sequence of statement forms. All statements in an argument and all statement forms in an argument form, except for the final one, are called **premises**. The final statement or statement form is called the **conclusion**.
- The symbol $\therefore$, which is read “therefore,” is normally placed just before the conclusion.
- To say that an **argument form** is **valid** means, **if the resulting premises are all true (AND), then the conclusion is also true.**

Determining Validity or Invalidity

- An argument form consisting of two premises and a conclusion is called **syllogism**.

If p then q

p

$\therefore q$

$(p \rightarrow q) \wedge p$

premises

p	q	$p \rightarrow q$	p	q
T	T	T	T	T
T	F	F	T	
F	T	T	F	
F	F	T	F	

Critical row

conclusion

It is a valid argument!

Rules of Inference

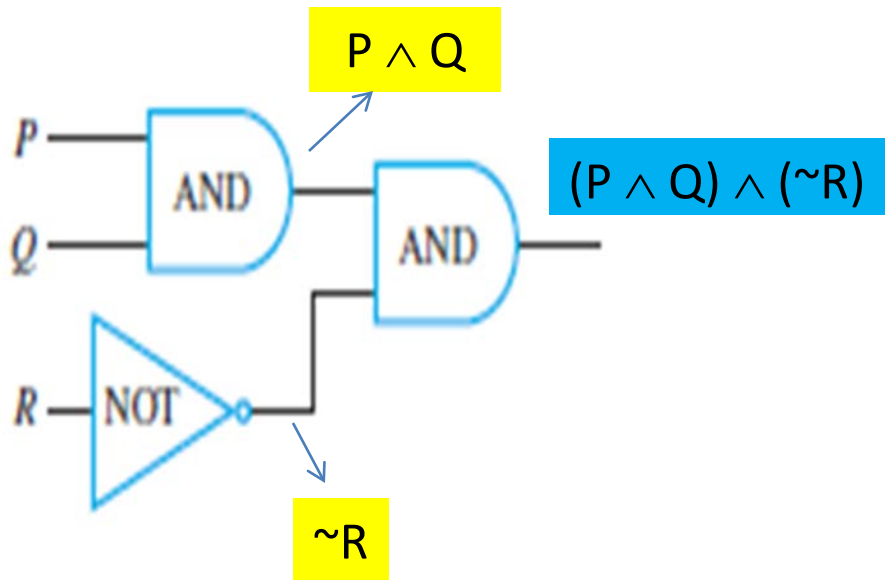
Modus Ponens	$p \rightarrow q$ p $\therefore q$	Elimination	a. $p \vee q$ $\sim q$ $\therefore p$ b. $p \vee q$ $\sim p$ $\therefore q$
Modus Tollens	$p \rightarrow q$ $\sim q$ $\therefore \sim p$	Transitivity	$p \rightarrow q$ $q \rightarrow r$ $\therefore p \rightarrow r$
Generalization	a. p $\therefore p \vee q$ b. q $\therefore p \vee q$	Proof by Division into Cases	$p \vee q$ $p \rightarrow r$ $q \rightarrow r$ $\therefore r$
Specialization	a. $p \wedge q$ $\therefore p$ b. $p \wedge q$ $\therefore q$		
Conjunction	p q $\therefore p \wedge q$	Contradiction Rule	$\sim p \rightarrow c$ $\therefore p$

Digital Logic Circuits

Type of Gate	Symbolic Representation	Action																		
NOT		<table><tr><th>Input</th><th>Output</th></tr><tr><th><i>P</i></th><th><i>R</i></th></tr><tr><td>1</td><td>0</td></tr><tr><td>0</td><td>1</td></tr></table>	Input	Output	<i>P</i>	<i>R</i>	1	0	0	1										
Input	Output																			
<i>P</i>	<i>R</i>																			
1	0																			
0	1																			
AND		<table><tr><th colspan="2">Input</th><th>Output</th></tr><tr><th><i>P</i></th><th><i>Q</i></th><th><i>R</i></th></tr><tr><td>1</td><td>1</td><td>1</td></tr><tr><td>1</td><td>0</td><td>0</td></tr><tr><td>0</td><td>1</td><td>0</td></tr><tr><td>0</td><td>0</td><td>0</td></tr></table>	Input		Output	<i>P</i>	<i>Q</i>	<i>R</i>	1	1	1	1	0	0	0	1	0	0	0	0
Input		Output																		
<i>P</i>	<i>Q</i>	<i>R</i>																		
1	1	1																		
1	0	0																		
0	1	0																		
0	0	0																		
OR		<table><tr><th colspan="2">Input</th><th>Output</th></tr><tr><th><i>P</i></th><th><i>Q</i></th><th><i>R</i></th></tr><tr><td>1</td><td>1</td><td>1</td></tr><tr><td>1</td><td>0</td><td>1</td></tr><tr><td>0</td><td>1</td><td>1</td></tr><tr><td>0</td><td>0</td><td>0</td></tr></table>	Input		Output	<i>P</i>	<i>Q</i>	<i>R</i>	1	1	1	1	0	1	0	1	1	0	0	0
Input		Output																		
<i>P</i>	<i>Q</i>	<i>R</i>																		
1	1	1																		
1	0	1																		
0	1	1																		
0	0	0																		

Boolean Expression for a Circuit

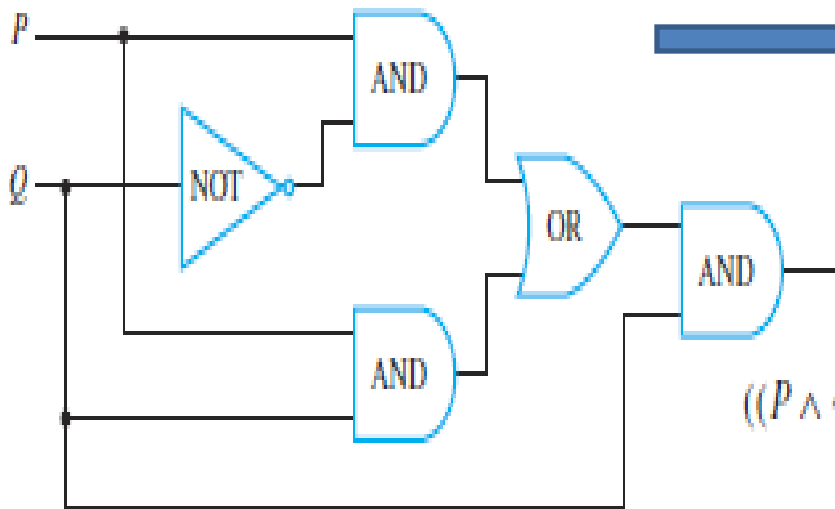
A statement variable or an input signal, that can take one of only two values is called a **Boolean variable**.



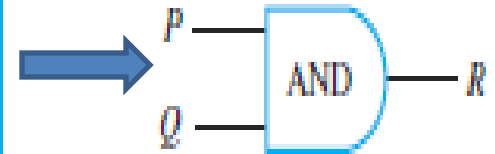
Input			Output
P	Q	R	V
1	1	1	0
1	1	0	1
1	0	1	0
1	0	0	0
0	1	1	0
0	1	0	0
0	0	1	0
0	0	0	0

Circuit for Boolean Expression

$$(P \wedge \sim Q) \vee (P \wedge Q) \wedge Q$$



Input		Output
P	Q	R
1	1	1
1	0	0
0	1	0
0	0	0



$((P \wedge \sim Q) \vee (P \wedge Q)) \wedge Q$ and $P \wedge Q$, respectively. By Theorem 2.1.1,

$$((P \wedge \sim Q) \vee (P \wedge Q)) \wedge Q$$

$$\equiv (P \wedge (\sim Q \vee Q)) \wedge Q \quad \text{by the distributive law}$$

$$\equiv (P \wedge (Q \vee \sim Q)) \wedge Q \quad \text{by the commutative law for } \vee$$

$$\equiv (P \wedge t) \wedge Q \quad \text{by the negation law}$$

$$\equiv P \wedge Q \quad \text{by the identity law.}$$

Showing that two circuits are equivalent