

Lecture 9: Division

Spring 2024
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Topics

- Dividing unsigned numbers
- Hardware designs for dividers
- Division in C

Unsigned Integer Division, Longhand

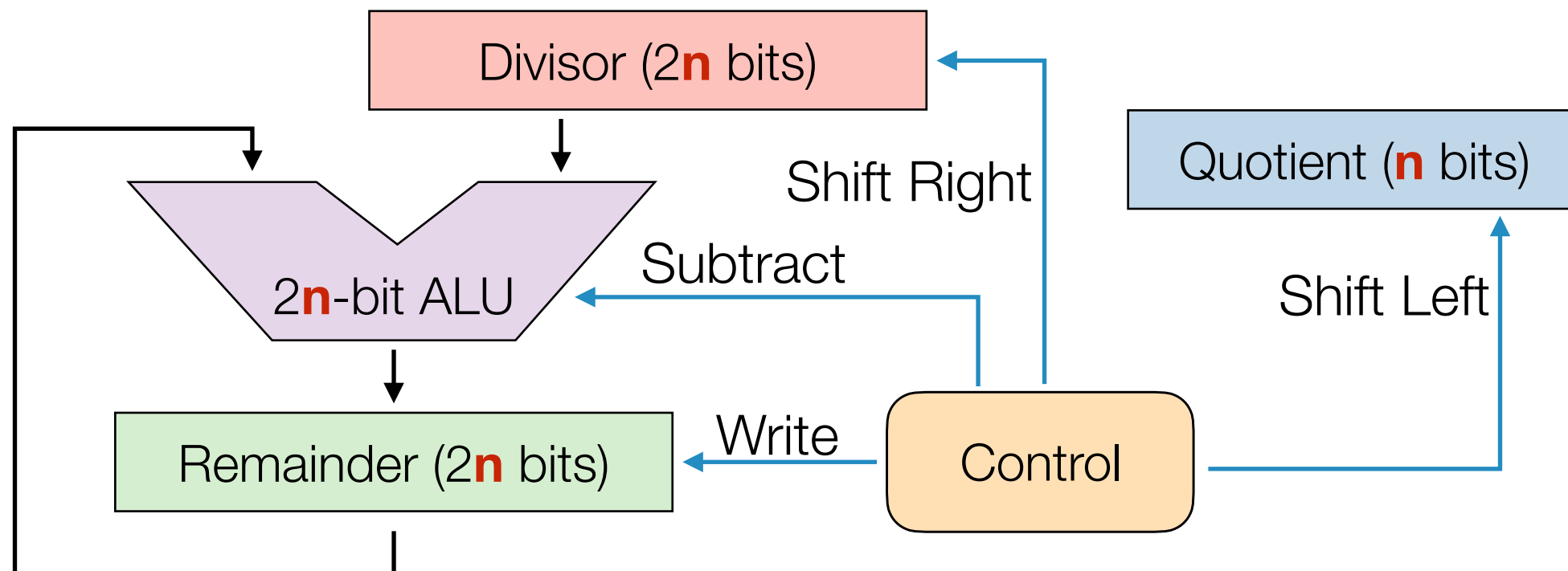
- Dividing binary values is like dividing decimals by hand:

$$\begin{array}{r}
 \begin{array}{c} = 8_{10} \\ \text{(Divisor)} \end{array}
 \begin{array}{c} 1\ 0\ 0\ 0 \end{array}
 \overline{) \begin{array}{c} \\ 1\ 0\ 0\ 1\ 0\ 1\ 0 \\ - 1\ 0\ 0\ 0 \\ \hline 1\ 0 \\ 1\ 0\ 1 \\ 1\ 0\ 1\ 0 \\ - 1\ 0\ 0\ 0 \\ \hline 1\ 0 \end{array}
 \begin{array}{c} = 9_{10} \text{ (Quotient)} \\ = 74_{10} \text{ (Dividend)} \\ = 2_{10} \text{ (Remainder or modulo result)} \end{array}
 \end{array}$$

- Dividend = Quotient $\times$ Divisor + Remainder
- Each bit in quotient means that divisor could be subtracted from partial dividend: 0 $\rightarrow$ smaller than divisor, 1 $\rightarrow$ greater than or equal to divisor

Unsigned Shift-Subtract Divider (version 1)

- Store **n**-bit divisor in a register twice its size, towards MSB
- For each step, shift divisor right and quotient left
- Initialize **2n**-bit remainder register to dividend, towards LSB
- **Control** decides when to shift and when to write new value into remainder register



Example of Shift-Subtract Divider (version 1)

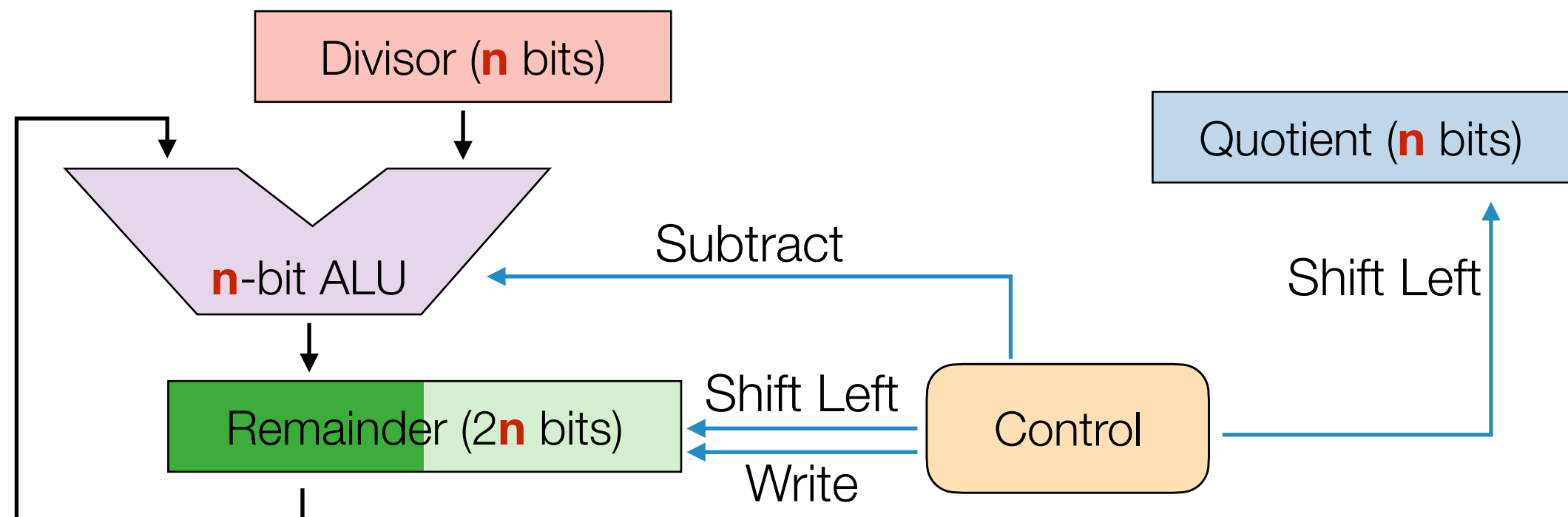
- Dividing two **n**-bit numbers needs **n+1** steps to generate **n**-bit quotient and remainder
 - Subtract divisor from remainder
 - If result is non-negative, keep new remainder and append 1 to quotient,
 - Else restore remainder, and append 0 to quotient
 - Logical shift divisor right, then shift quotient left
 - Repeat **n+1** times

Example: $7_{10} \div 2_{10}$, 4-bits

Iteration	Operations	Quotient	Remainder	Divisor
0	Initial	0000	0000 0111	0010 0000
1a	Subtract	0000	1110 0111	0010 0000
1b	Restore remainder	0000	0000 0111	0010 0000
1c	Logical shift divisor right	0000	0000 0111	0001 0000
1d	Shift quotient left	0000	0000 0111	0001 0000
after 2	Sub, restore, shift D, shift Q	0000	0000 0111	0000 1000
after 3	Sub, restore, shift D, shift Q	0000	0000 0111	0000 0100
after 4	Sub, keep, shift D, shift Q	0001	0000 0011	0000 0010
after 5	Sub, keep, shift D, shift Q	0011	0000 0001	0000 0001

Unsigned Shift-Subtract Divider (version 2)

- Because half of divisor register is filled with zeroes, $2n$ -bit ALU is wasteful
- Use only n -bit divisor, n -bit ALU; keep $2n$ -bit remainder
- Result of ALU is written to **upper half** of remainder register
- After each step, control **shifts left** the remainder register



Example of Shift-Subtract Divider (version 2)

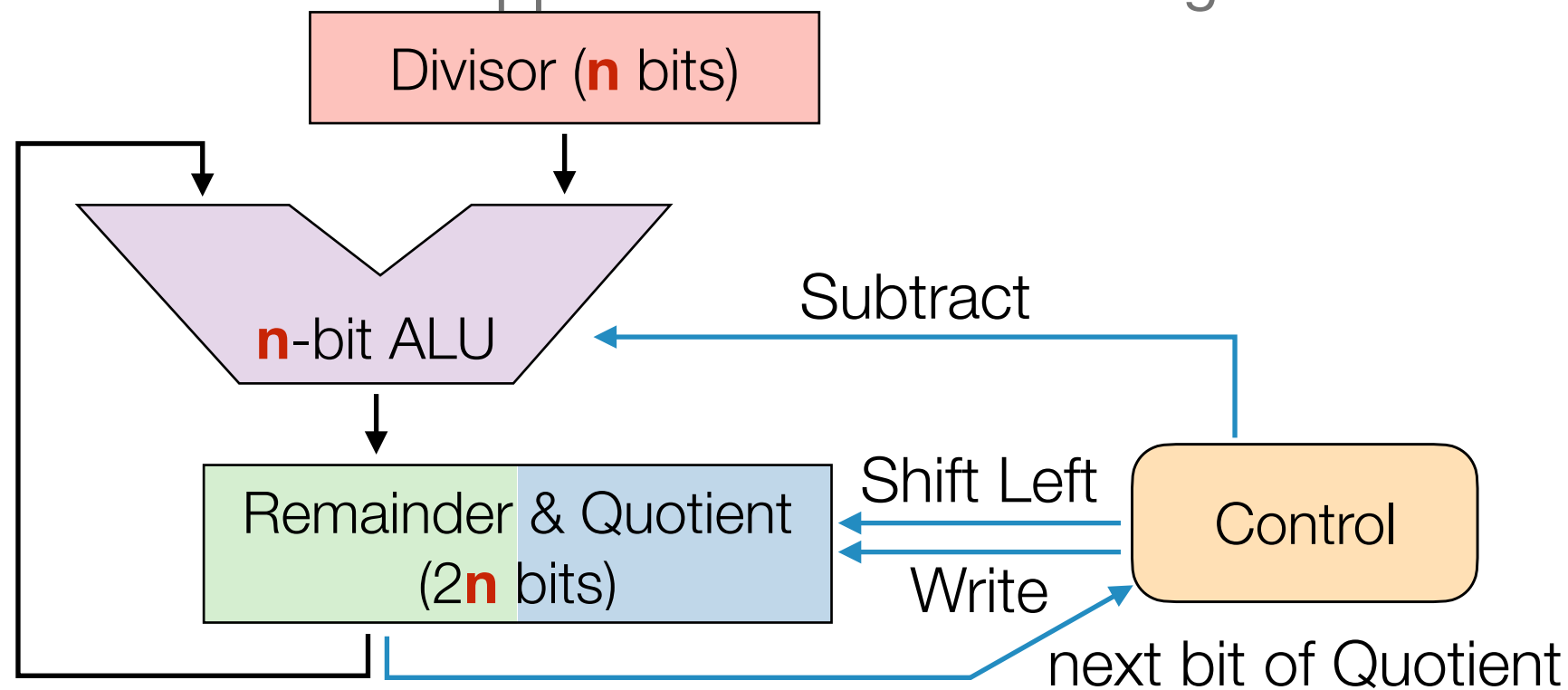
- Subtract from upper half of remainder
- If test subtraction is non-negative, keep result in upper half of remainder
- Because iteration 1 cannot produce a 1 in quotient, first shift remainder and then subtract to save an iteration
- Repeat **n** (instead of **n**+1) times

Example: $7_{10} \div 2_{10}$, 4-bits

Iteration	Operations	Quotient	Remainder	Divisor
0	Initial	0000	0000 0111	0010
1a	Shift remainder left	0000	0000 1110	0010
1b	Subtract	0000	1110 1110	0010
1c	Restore	0000	0000 1110	0010
1d	Shift quotient left	0000	0000 1110	0010
after 2	Shift R, sub, restore, shift Q	0000	0001 1100	0010
after 3	Shift R, sub, keep, shift Q	0001	0001 1000	0010
after 4	Shift R, sub, keep, shift Q	0011	0001 0000	0010

Unsigned Shift-Subtract Divider (version 3)

- At startup, top half of remainder register will be shifted away and is unused
- Combine quotient and remainder registers, with remainder in top half of register and quotient in bottom half
- ALU results still written to upper half of combined register



Example of Shift-Subtract Divider (version 3)

- After **n** iterations, upper half of combined register holds the remainder, bottom half holds quotient

Example: $7_{10} \div 2_{10}$, 4-bits

Iteration	Operations	Remainder / Quotient	Divisor
0	Initial	0000 0111	0010
1a	Shift R/Q left	0000 1110	0010
1b	Subtract	1110 1110	0010
1c	Restore	0000 1110	0010
after 2	Shift R/Q, sub, restore	0001 1100	0010
after 3	Shift R/Q, sub, keep	0001 1001	0010
after 4	Shift R/Q, sub, keep	0001 0011	0010

Multiplication versus Division

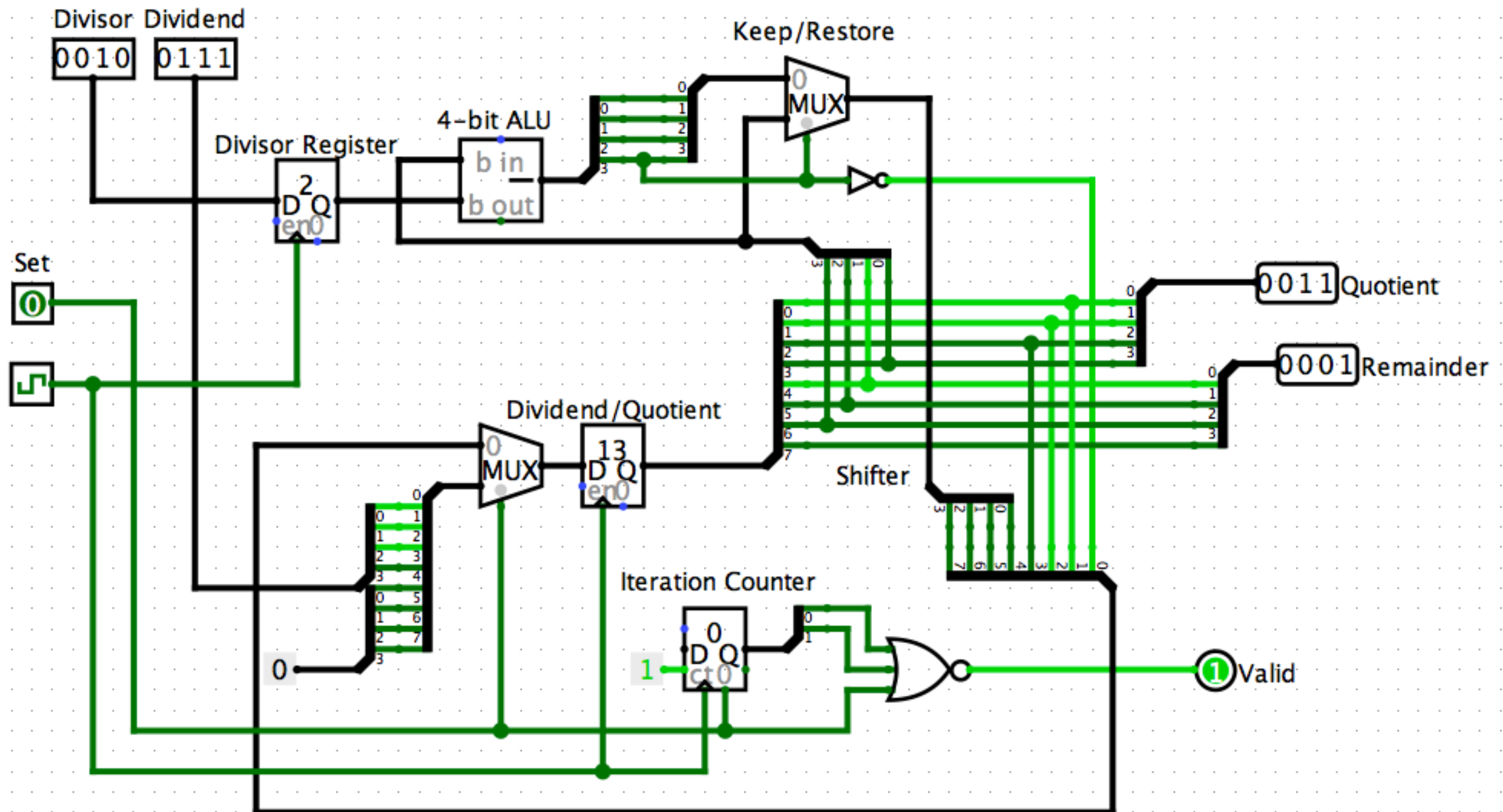
- Hardware multipliers can be optimized to run in only a few cycles
 - Parallel adders; algorithm tricks like *Booth*, *Karatsuba*, and *Toom-Cook*
- Hardware dividers all rely upon some kind of iterative approach
- For a typical ARMv8-A implementation:

Instruction Type	Number of Cycles
Integer add or subtract	1
Integer multiply	3
Multiply-accumulate	4
Integer division	12

Optimizing Dividers

- Instead of subtracting the divisor, add the divisor's inverse
- Instead of keeping/restoring the combined register, use a mux to select which bits to write back to the register
 - If MSB out of ALU is 0, then result is non-negative, which means to keep change
 - Else if MSB is 1, then result is negative, which is a restore
- Some implementations will check if the divisor is zero
 - Will raise an [interrupt](#) if a dividing by zero

Unsigned Shift-Subtract Divider Implementation



Signed Division

- Simplest approach is to remember signs, make both dividend and divisor positive, and then apply correct signs afterwards

- Quotient's sign is negative when dividend and divisor's sign differ

- $Q_{\text{sign}} = \text{Dividend}_{\text{sign}} \oplus \text{Divisor}_{\text{sign}}$

- Remainder's sign is same as dividend's sign

Dividend		Divisor		Result
7	÷	2	=	+3, remainder +1
7	÷	-2	=	-3, remainder +1
-7	÷	2	=	-3, remainder -1
-7	÷	-2	=	+3, remainder -1

- Otherwise, $-7 \div 2$ could be -4, remainder +1

Division in C

- Result of `/` operator has the same sign as the dividend
 - In C99, `/` rounds towards 0 (relevant with a negative operand)
- The `%` operator is the **modulo** operator
- By definition, $(a / b) * b + (a \% b) = a$

```
void f(int a, int b) {  
    printf("%d / %d = %d, mod %d\n",  
          a, b, (a/b), (a%b));  
}  
int main(void) {  
    f(7, 2);      f(7, -2);  
    f(-7, 2);     f(-7, -2);  
    return 0;  
}
```

7	/	2	=	3	, mod 1
7	/	-2	=	-3	, mod 1
-7	/	2	=	-3	, mod -1
-7	/	-2	=	3	, mod -1

Division Instructions

- Two different types of division:

Division Type	ARMv8-A	PowerPC	x86-64
Signed Division	<code>sdiv</code>	<code>divw</code>	<code>idiv</code>
Unsigned Division	<code>udiv</code>	<code>divwu</code>	<code>div</code>

- For ARMv8-A and PowerPC, **no** hardware check for divide by zero
- For ARMv8-A and PowerPC, **no** builtin way to get remainder / modulo

More Detailed Division Code Example

```
#include <stdio.h>
#include <stdlib.h>

int main(int argc, char *argv[]) {
    signed long s1 = strtol(argv[1], NULL, 0);
    signed long s2 = strtol(argv[2], NULL, 0);
    unsigned long u1 = strtoul(argv[1], NULL, 0);
    unsigned long u2 = strtoul(argv[2], NULL, 0);

    signed long squot = s1 / s2;
    signed long smod = s1 % s2;
    unsigned long uquot = u1 / u2;
    unsigned long umod = u1 % u2;
    printf("    signed a / b = %ld, a %% b = %ld\n", squot, smod);
    printf("unsigned a / b = %lu, a %% b = %lu\n", uquot, umod);
    return 0;
}
```

x86-64 Division

- On x86-64, 128-bit signed division is performed by the **idivq** instruction:
 - **rdx** holds upper 64 bits of dividend
 - **rax** holds lower 64 bits of dividend
 - Operand to **idivq** is the divisor
- Afterwards, quotient will be in **rax**, modulo in **rdx**

x86-64 Division

```
    signed long squot = s1 / s2;
400649:      48 8b 45 c0      mov     -0x40(%rbp),%rax
40064d:      48 99            cqto
40064f:      48 f7 7d c8      idivq   -0x38(%rbp)
400653:      48 89 45 e0      mov     %rax,-0x20(%rbp)

    signed long smod = s1 % s2;
400657:      48 8b 45 c0      mov     -0x40(%rbp),%rax
40065b:      48 99            cqto
40065d:      48 f7 7d c8      idivq   -0x38(%rbp)
400661:      48 89 55 e8      mov     %rdx,-0x18(%rbp)

    unsigned long uquot = u1 / u2;
400665:      48 8b 45 d0      mov     -0x30(%rbp),%rax
400669:      ba 00 00 00 00    mov     $0x0,%edx
40066e:      48 f7 75 d8      divq    -0x28(%rbp)
400672:      48 89 45 f0      mov     %rax,-0x10(%rbp)

    unsigned long umod = u1 % u2;
400676:      48 8b 45 d0      mov     -0x30(%rbp),%rax
40067a:      ba 00 00 00 00    mov     $0x0,%edx
40067f:      48 f7 75 d8      divq    -0x28(%rbp)
400683:      48 89 55 f8      mov     %rdx,-0x8(%rbp)
```

Intel Division Implementation

- Original implementation used a shift-subtract divider, about half the speed of shift-adder multiplication
- Newer chips use a **SRT table** to look up a quotient digit, based upon dividend and divisor values
 - Intel Pentium chips had a famous hardware bug due to incorrectly encoded SRT table

ARMv8-A Division

- On ARMv8-A, 64-bit signed division is performed by the **sdiv** instruction
 - Divisor, dividend, and quotient registers are all specified by instruction operands
- Upon a divide-by-zero, 0 is written to quotient register, with no other indication that any problem occurred
- To get the modulo, multiply the quotient by dividend, and then subtract that from the divisor

ARMv8-A Division

```
    signed long squot = s1 / s2;
4006e0:    f94013a1    ldr    x1, [x29, #32]
4006e4:    f94017a0    ldr    x0, [x29, #40]
4006e8:    9ac00c20    sdiv   x0, x1, x0
4006ec:    f90023a0    str    x0, [x29, #64]
    signed long smod = s1 % s2;
4006f0:    f94013a0    ldr    x0, [x29, #32]
4006f4:    f94017a1    ldr    x1, [x29, #40]
4006f8:    9ac10c02    sdiv   x2, x0, x1
4006fc:    f94017a1    ldr    x1, [x29, #40]
400700:    9b017c41    mul    x1, x2, x1
400704:    cb010000    sub    x0, x0, x1
400708:    f90027a0    str    x0, [x29, #72]
    unsigned long uquot = u1 / u2;
40070c:    f9401ba1    ldr    x1, [x29, #48]
400710:    f9401fa0    ldr    x0, [x29, #56]
400714:    9ac00820    udiv   x0, x1, x0
400718:    f9002ba0    str    x0, [x29, #80]
```